# Computational Intelligence 

## Winter Term 2014/15

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## Plan for Today

- Bidirectional Associative Memory (BAM)
- Fixed Points
- Concept of Energy Function
- Stable States = Minimizers of Energy Function
- Hopfield Network
- Convergence
- Application to Combinatorial Optimization


## Bidirectional Associative Memory (BAM)

## Network Model



- fully connected
- bidirectional edges
- synchonized:
step t : data flow from x to y
step $\mathrm{t}+1$ : data flow from y to x

$$
\text { start: } \begin{aligned}
y^{(0)} & =\operatorname{sgn}\left(x^{(0)} W\right) \\
& x^{(1)}=\operatorname{sgn}\left(y^{(0)} W\right) \\
& y^{(1)}=\operatorname{sgn}\left(x^{(1)} W\right) \\
& x^{(2)}=\operatorname{sgn}\left(y^{(1)} W\right)
\end{aligned}
$$

bipolar inputs $\in\{-1,+1\}$

## Bidirectional Associative Memory (BAM)

## Fixed Points

## Definition

$(x, y)$ is fixed point of BAM iff $y=\operatorname{sgn}(x W)$ and $x^{\prime}=\operatorname{sgn}\left(W y^{\prime}\right)$.

Set $\mathrm{W}=\mathrm{x}^{\prime} \mathrm{y} . \quad$ (note: x is row vector)
$y=\operatorname{sgn}(x W)=\operatorname{sgn}\left(x\left(x^{\prime} y\right)\right)=\operatorname{sgn}\left(\left(x x^{\prime}\right) y\right)=\operatorname{sgn}(\underbrace{\|x\|^{2} y})=y$
$>0$ (does not alter sign)
$x^{\prime}=\operatorname{sgn}\left(W y^{\prime}\right)=\operatorname{sgn}\left(\left(x^{\prime} y\right) y^{\prime}\right)=\operatorname{sgn}\left(x^{x^{\prime}}\left(y y^{\prime}\right)\right)=\operatorname{sgn}(x^{x^{\prime}} \underbrace{\|y\|^{2}}_{>0 \text { (does not alter sign) }})=x^{\prime}$
Theorem: If $W=x^{\prime} y$ then $(x, y)$ is fixed point of BAM.

## Bidirectional Associative Memory (BAM)

## Concept of Energy Function

given: $B A M$ with $W=x^{\prime} y \quad \Rightarrow(x, y)$ is stable state of BAM
starting point $\mathrm{x}^{(0)}$

$$
\begin{aligned}
& \Rightarrow \mathrm{y}^{(0)}=\operatorname{sgn}\left(\mathrm{x}^{(0)} \mathrm{W}\right) \\
& \Rightarrow \text { excitation } \mathrm{e}^{‘}=\mathrm{W}\left(\mathrm{y}^{(0)}\right)^{،}
\end{aligned}
$$

$\Rightarrow$ if $\operatorname{sign}\left(e^{d}\right)=x^{(0)}$ then $\left(x^{(0)}, y^{(0)}\right)$ stable state

recall: $\frac{a b^{\prime}}{\|a\| \cdot\|b\|}=\cos \angle(a, b)$

small angle $\alpha \Rightarrow$ large $\cos (\alpha)$

## Bidirectional Associative Memory (BAM)

## Concept of Energy Function

required:
small angle between $e^{‘}=W y^{(0)}$ ‘ and $x^{(0)}$
$\Rightarrow$ larger cosine of angle indicates greater similarity of vectors
$\Rightarrow \forall \mathrm{e}^{‘}$ of equal size: try to maximize $\mathrm{x}^{(0)} \mathrm{e}^{‘}=\left\|\mathrm{x}^{(0)}\right\| \cdot\|\mathrm{e}\| \cdot \cos \angle\left(\mathrm{x}^{(0)}, \mathrm{e}\right)$
fixed fixed $\rightarrow$ max!
$\Rightarrow$ maximize $x^{(0)} e^{‘}=x^{(0)} W y^{(0)}{ }^{\text {‘ }}$
$\Rightarrow$ identical to minimize $-x^{(0)} \mathrm{W} \mathrm{y}^{(0)}$ ‘

## Definition

Energy function of BAM at iteration $t$ is $E\left(x^{(t)}, y^{(t)}\right)=-\frac{1}{2} x^{(t)} W y^{(t)}$ ‘

## Bidirectional Associative Memory (BAM)

Stable States

## Theorem

An asynchronous BAM with arbitrary weight matrix W reaches steady state in a finite number of updates.

Proof:
$E(x, y)=-\frac{1}{2} x W y^{\prime}=\left\{\begin{aligned} &-\frac{1}{2} x\left(W y^{\prime}\right)=-\frac{1}{2} x b^{\prime}=-\frac{1}{2} \sum_{i=1}^{n} b_{i} x_{i} \\ &-\frac{1}{2}(x W) y^{\prime}=-\frac{1}{2} a y^{\prime} \\ &=-\frac{1}{2} \sum_{i=1}^{k} a_{i} y_{i}\end{aligned}\right.$ excitations

BAM asynchronous $\Rightarrow \quad$ select neuron at random from left or right layer, compute its excitation and change state if necessary (states of other neurons not affected)

## Bidirectional Associative Memory (BAM)

$$
\text { neuron } \begin{aligned}
i \text { of left layer has changed } & \Rightarrow \operatorname{sgn}\left(x_{i}\right) \neq \operatorname{sgn}\left(b_{i}\right) \\
& \Rightarrow x_{i} \text { was updated to } \widetilde{x}_{i}=-x_{i}
\end{aligned}
$$

$$
E(x, y)-E(\tilde{x}, y)=-\frac{1}{2} \underbrace{b_{i}\left(x_{i}-\tilde{x}_{i}\right)}_{<0}>0
$$

| $x_{i}$ | $b_{i}$ | $x_{i}-\widetilde{x}_{i}$ |
| :---: | :---: | :---: |
| -1 | $>0$ | $<0$ |
| +1 | $<0$ | $>0$ |

use analogous argumentation if neuron of right layer has changed
$\Rightarrow$ every update (change of state) decreases energy function
$\Rightarrow$ since number of different bipolar vectors is finite update stops after finite \#updates
remark: dynamics of BAM get stable in local minimum of energy function!

## Hopfield Network

special case of BAM but proposed earlier (1982)

## characterization:

- neurons preserve state until selected at random for update

- n neurons fully connected
- symmetric weight matrix
- no self-loops ( $\rightarrow$ zero main diagonal entries)
- thresholds $\theta$, neuron i fires if excitations larger than $\theta_{\mathrm{i}}$

transition: select index k at random, new state is $\tilde{x}=\operatorname{sgn}(x W-\theta)$

$$
\text { where } \tilde{x}=\left(x_{1}, \ldots, x_{k-1}, \tilde{x}_{k}, x_{k+1}, \ldots, x_{n}\right)
$$

energy of state x is $E(x)=-\frac{1}{2} x W x^{\prime}+\theta x^{\prime}$

## Hopfield Network

## Lecture 04

## Theorem:

Hopfield network converges to local minimum of energy function after a finite number of updates.

Proof: $\quad$ assume that $\mathrm{x}_{\mathrm{k}}$ has been updated $\Rightarrow \tilde{x}_{k}=-x_{k}$ and $\tilde{x}_{i}=x_{i}$ for $i \neq k$
$E(x)-E(\tilde{x})=-\frac{1}{2} x W x^{\prime}+\theta x^{\prime}+\frac{1}{2} \tilde{x} W \tilde{x}^{\prime}-\theta \tilde{x}^{\prime}$
$=-\frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j} x_{i} x_{j}+\sum_{i=1}^{n} \theta_{i} x_{i}+\frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j} \tilde{x}_{i} \tilde{x}_{j}-\sum_{i=1}^{n} \theta_{i} \tilde{x}_{i}$
$=-\frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j}\left(x_{i} x_{j}-\tilde{x}_{i} \tilde{x}_{j}\right)+\sum_{i=1}^{n} \theta_{i} \underbrace{\left(x_{i}-\tilde{x}_{i}\right)}_{=0 \text { if } \mathrm{i} \neq \mathrm{k}}$
$=-\frac{1}{2} \sum_{\substack{i=1 \\ i \neq k}}^{n} \sum_{\substack{ \\\|_{i}}}^{n} w_{i j}\left(x_{i} x_{j}-\tilde{x}_{i} \tilde{x}_{j}\right)-\frac{1}{2} \sum_{j=1}^{n} w_{\substack{ \\x_{i}}}^{n}\left(x_{k} x_{j}-\tilde{x}_{k} \tilde{x}_{j}\right)+\theta_{k}\left(x_{k}-\tilde{x}_{k}\right)$

$$
>0 \text { if } x_{k}<0 \text { and vice versa }
$$

$$
\begin{aligned}
& =-\frac{1}{2} \sum_{\substack{i=1 \\
i \neq k}}^{n} \sum_{j=1}^{n} w_{i j} x_{i} \underbrace{\left(x_{j}-\tilde{x}_{j}\right)}_{=0 \text { if } \mathrm{j} \neq \mathrm{k}}-\frac{1}{2} \sum_{\substack{j=1 \\
j \neq k}}^{n} w_{k j} x_{j}\left(x_{k}-\tilde{x}_{k}\right)+\theta_{k}\left(x_{k}-\tilde{x}_{k}\right) \\
& =-\frac{1}{2} \sum_{\substack{i=1 \\
i \neq k}}^{n} w_{i k} x_{i}\left(x_{k}-\tilde{x}_{k}\right)-\frac{1}{2} \sum_{\substack{j=1 \\
j \neq k}}^{n} w_{k j} x_{j}\left(x_{k}-\tilde{x}_{k}\right)+\theta_{k}\left(x_{k}-\tilde{x}_{k}\right) \\
& =-\sum_{i=1}^{n} w_{i k} x_{i}\left(x_{k}-\tilde{x}_{k}\right)+\theta_{k}\left(x_{k}-\tilde{x}_{k}\right) \\
& =-\left(x_{k}-\tilde{x}_{k}\right)[\underbrace{\left[\sum_{i=1}^{n} w_{i k} x_{i}\right.}_{\text {excitation } \mathrm{e}_{\mathrm{k}}}-\theta_{k}]>0 \quad \begin{array}{l}
\text { since: } \\
\begin{array}{cccc}
x_{k} & x_{k}-\tilde{x}_{k} & e_{k}-\theta_{k} & \Delta E
\end{array} \\
\begin{array}{cccccc}
+1 & >0 & <0 & >0 \\
-1 & <0 & >0 & >0
\end{array}
\end{array}
\end{aligned}
$$

## Hopfield Network

## Application to Combinatorial Optimization

## Idea:

- transform combinatorial optimization problem as objective function with $x \in\{-1,+1\}^{n}$
- rearrange objective function to look like a Hopfield energy function
- extract weights W and thresholds $\theta$ from this energy function
- initialize a Hopfield net with these parameters W and $\theta$
- run the Hopfield net until reaching stable state (= local minimizer of energy function)
- stable state is local minimizer of combinatorial optimization problem


## Hopfield Network

## Example I: Linear Functions

$$
f(x)=\sum_{i=1}^{n} c_{i} x_{i} \quad \rightarrow \min !\quad\left(x_{i} \in\{-1,+1\}\right)
$$

Evidently: $E(x)=f(x)$ with $W=0$ and $\theta=c$
$\Downarrow$
choose $x^{(0)} \in\{-1,+1\}^{n}$
set iteration counter $t=0$
repeat
choose index $k$ at random

$$
x_{k}^{(t+1)}=\operatorname{sgn}\left(x^{(t)} \cdot W_{\cdot, k}-\theta_{k}\right)=\operatorname{sgn}\left(x^{(t)} \cdot 0-c_{k}\right)=-\operatorname{sgn}\left(c_{k}\right)=\left\{\begin{aligned}
-1 & \text { if } c_{k}>0 \\
+1 & \text { if } c_{k}<0
\end{aligned}\right.
$$

increment $t$
until reaching fixed point
$\Rightarrow$ fixed point reached after $\Theta(\mathrm{n} \log \mathrm{n})$ iterations on average

## Hopfield Network

## Example II: MAXCUT

given: graph with $n$ nodes and symmetric weights $\omega_{\mathrm{ij}}=\omega_{\mathrm{ji}}$, $\omega_{\mathrm{ii}}=0$, on edges
task: find a partition $\mathrm{V}=\left(\mathrm{V}_{0}, \mathrm{~V}_{1}\right)$ of the nodes such that the weighted sum of edges with one endpoint in $\mathrm{V}_{0}$ and one endpoint in $\mathrm{V}_{1}$ becomes maximal

$$
\text { encoding: } \forall \mathrm{i}=1, \ldots, \mathrm{n}: \quad \mathrm{y}_{\mathrm{i}}=0 \Leftrightarrow \text { node } \mathrm{i} \text { in set } \mathrm{V}_{0} ; \quad \mathrm{y}_{\mathrm{i}}=1 \Leftrightarrow \text { node } \mathrm{i} \text { in set } \mathrm{V}_{1}
$$

objective function: $f(y)=\sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j}\left[y_{i}\left(1-y_{j}\right)+y_{j}\left(1-y_{i}\right)\right] \quad \rightarrow \max !$

## preparations for applying Hopfield network

step 1: conversion to minimization problem
step 2: transformation of variables
step 3: transformation to "Hopfield normal form"
step 4: extract coefficients as weights and thresholds of Hopfield net

## Hopfield Network

## Example II: MAXCUT (continued)

step 1: conversion to minimization problem

$$
\Rightarrow \text { multiply function with }-1 \quad \Rightarrow \mathrm{E}(\mathrm{y})=-\mathrm{f}(\mathrm{y}) \quad \rightarrow \text { min! }
$$

step 2: transformation of variables

$$
\Rightarrow y_{i}=\left(x_{i}+1\right) / 2
$$

$$
\Rightarrow f(x)=\sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j}\left[\frac{x_{i}+1}{2}\left(1-\frac{x_{j}+1}{2}\right)+\frac{x_{j}+1}{2}\left(1-\frac{x_{i}+1}{2}\right)\right]
$$

$$
=\frac{1}{2} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j}\left[1-x_{i} x_{j}\right]
$$

$$
=\underbrace{\frac{1}{2} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j}}-\frac{1}{2} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j} x_{i} x_{j}
$$

constant value (does not affect location of optimal solution)

## Hopfield Network

## Example II: MAXCUT (continued)

step 3: transformation to "Hopfield normal form"

$$
\begin{aligned}
E(x) & =\frac{1}{2} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \omega_{i j} x_{i} x_{j}=-\frac{1}{2} \sum_{i=1}^{n} \sum_{i \neq 1}^{n} \underbrace{\left(-\frac{1}{2} \omega_{i j}\right.}_{\mathrm{w}_{\mathrm{ij}}}) x_{i} x_{j} \\
& =-\frac{1}{2} x^{\prime} W x+\theta^{\prime} x
\end{aligned}
$$

step 4: extract coefficients as weights and thresholds of Hopfield net
$w_{i j}=-\frac{\omega_{i j}}{2}$ for $i \neq j, \quad w_{i i}=0, \quad \theta_{i}=0$
remark: $\omega_{i j}$ : weights in graph - $w_{i j}:$ weights in Hopfield net

